

# XS-Stabilizer Formalism

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joint work with Buerschaper, Van den Nest

# Outline

- Motivation
- Example: double semion model
- Summary of properties

# Definition

- Pauli-S group:  $\mathcal{P}_n^s = \langle \alpha, X, S \rangle^{\otimes n}$

$$\alpha = \sqrt{i} \quad S = \text{diag}(1, i) \quad S^{-1}XS = -iXZ$$
$$X \otimes S \otimes Z$$

- Given  $G = \langle g_1, \dots, g_m \rangle \subset \mathcal{P}_n^s$

We call a state  $|\psi\rangle$  XS-stabilizer state if (uniquely)

$$g_j|\psi\rangle = |\psi\rangle$$

When not unique, we call it XS-stabilizer code

# Motivation

# Pauli stabilizer formalism

- (Innocently looking) tensor product operators
- Most properties from commutation relation and linear algebra
- Numerous applications: Fault tolerance, measurement based computation, etc

# XS stabilizer

- (Still innocently looking) tensor product operators
- Many properties from commutation relation and linear algebra
- Simple to learn

# Toric (surface) code

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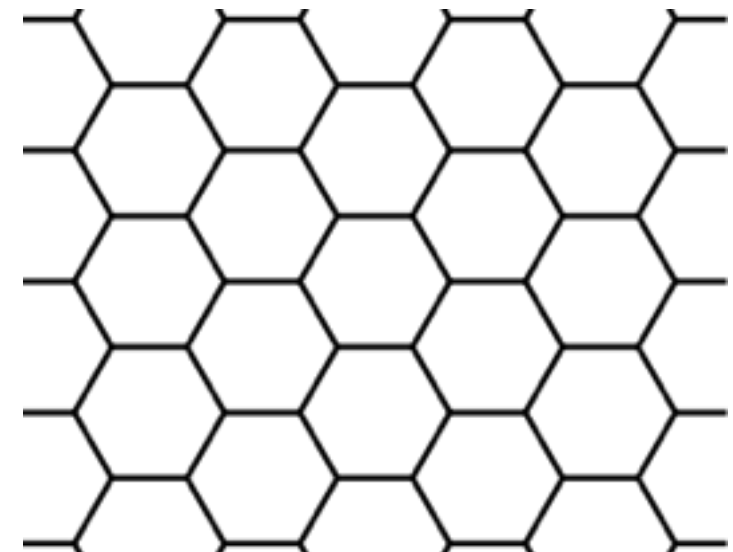
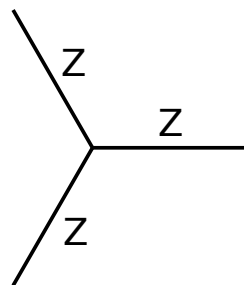
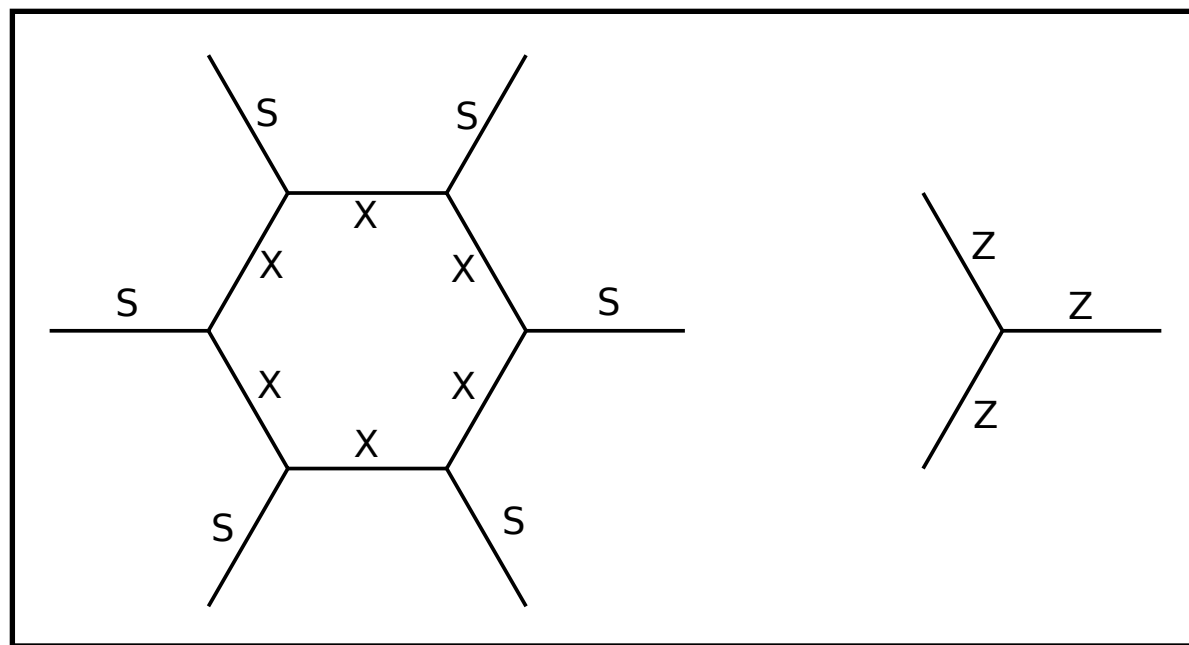
- Practical way to build an active quantum memory



# Toric (surface) code

- Practical way to build an active quantum memory
- Great example to understand basic properties of systems with topological order
- Exactly solvable and simple
- Contains features like anyons, string operators, boundary, twist, etc.

# XS-stabilizer: double semion and more



# Other motivations

# Other motivations

- (Efficient) representation of a larger class of quantum states

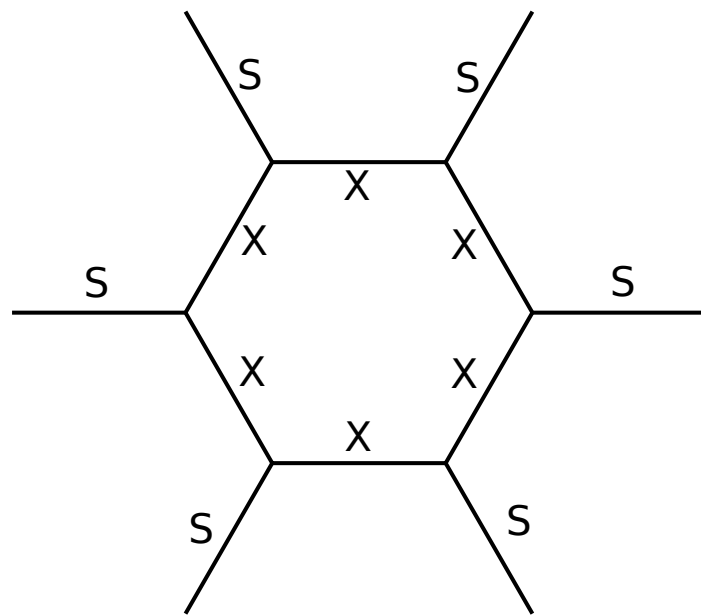
# Other motivations

- (Efficient) representation of a larger class of quantum states
- A class of commuting projector problems that are in **NP** (**P**)

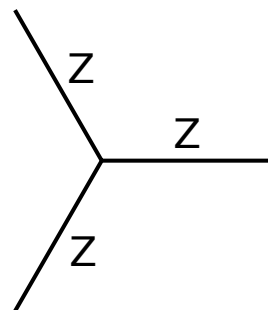
# An introduction to the Double semion model

# Double semion model

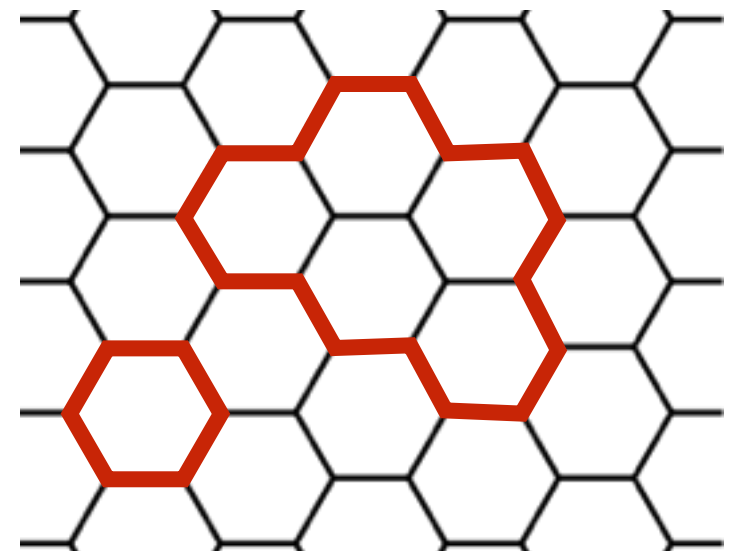
$$\sum_{x \text{ is close loops}} (-1)^{\text{number of loops}} |x\rangle$$



$g_p$



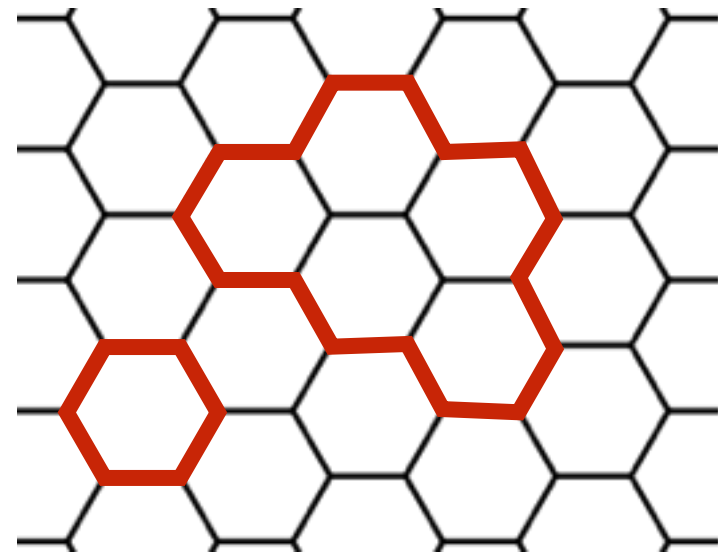
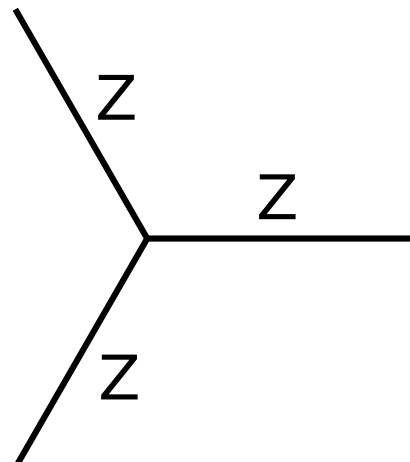
$g_v$



$|0\rangle$   $|1\rangle$

$$g_p |\psi\rangle = g_v |\psi\rangle = |\psi\rangle$$

# Z-type operator



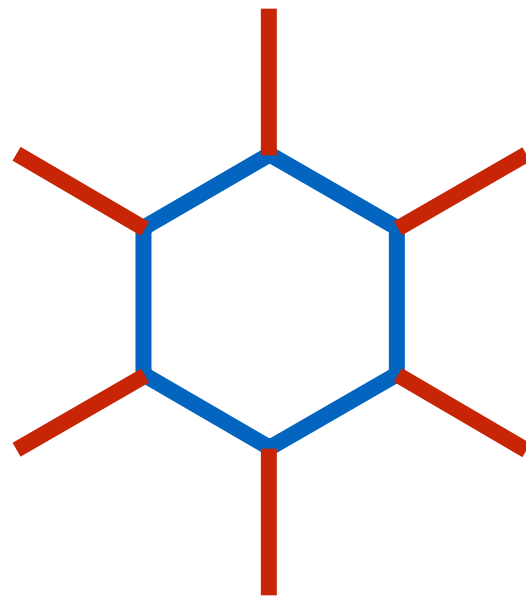
$|0\rangle$   $|1\rangle$

Gauge invariant subspace

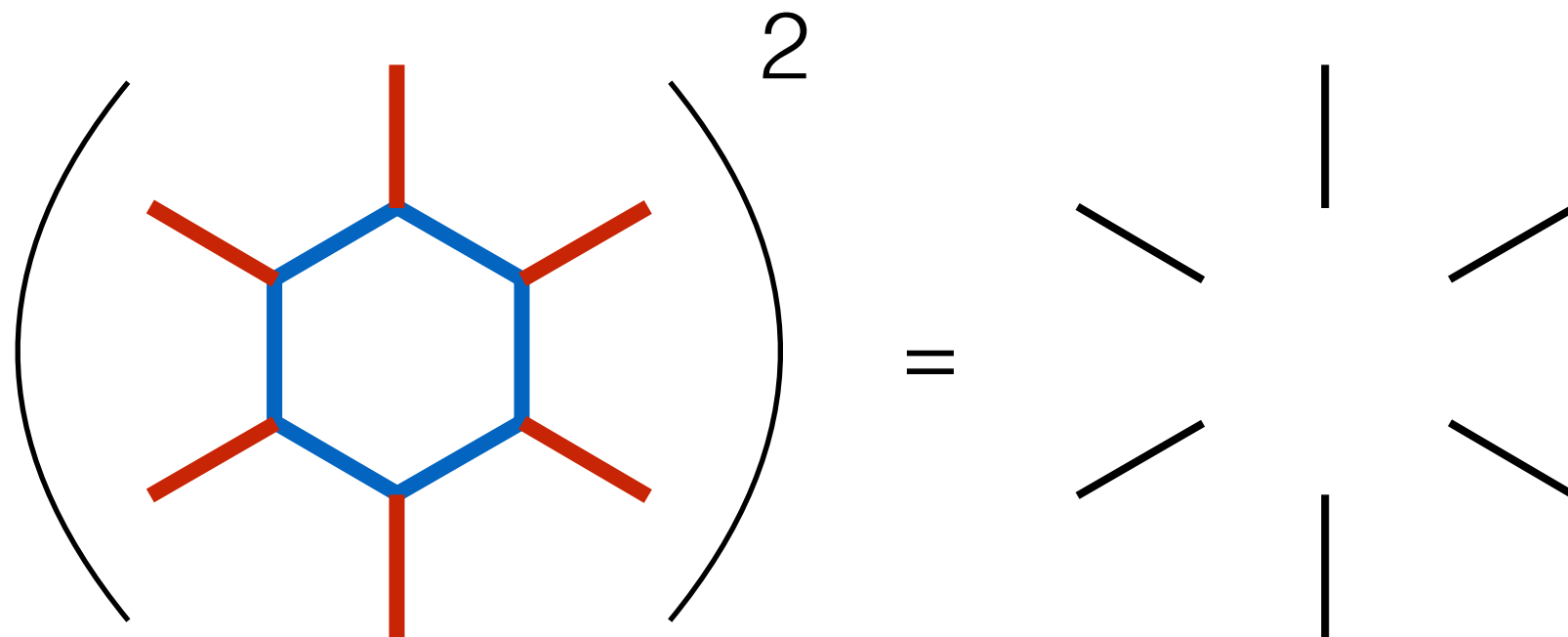


# Plaquette operator

$x$   
 $s$   
 $z$

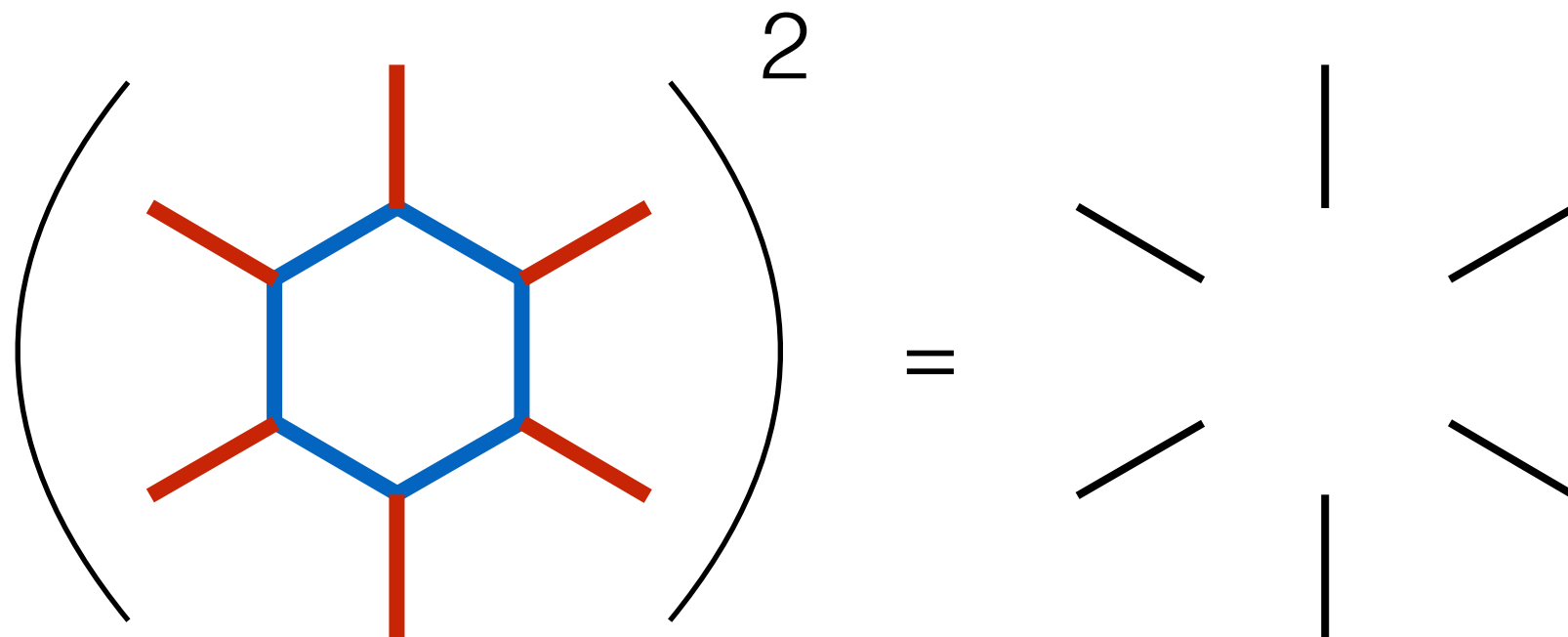


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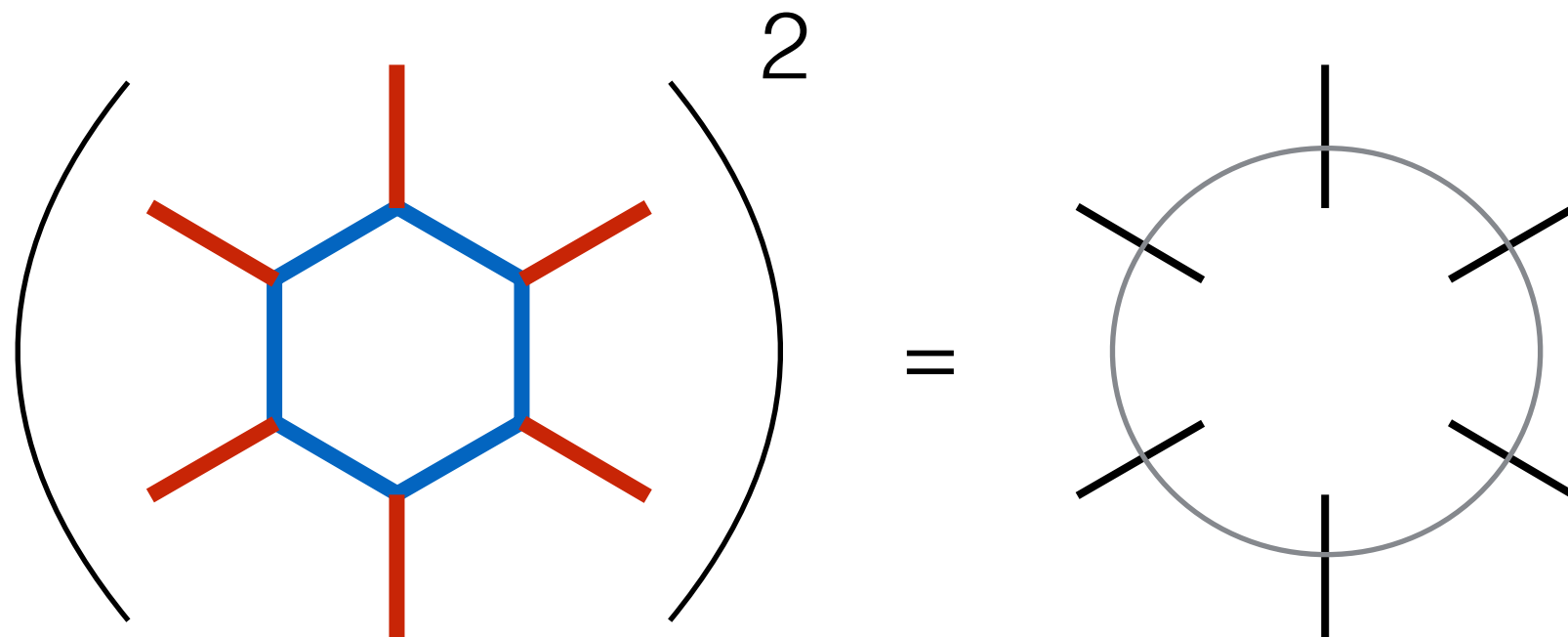


**X**  
**S**  
**Z**

- The square is equal to 1 inside gauge invariant subspace

# Plaquette operator

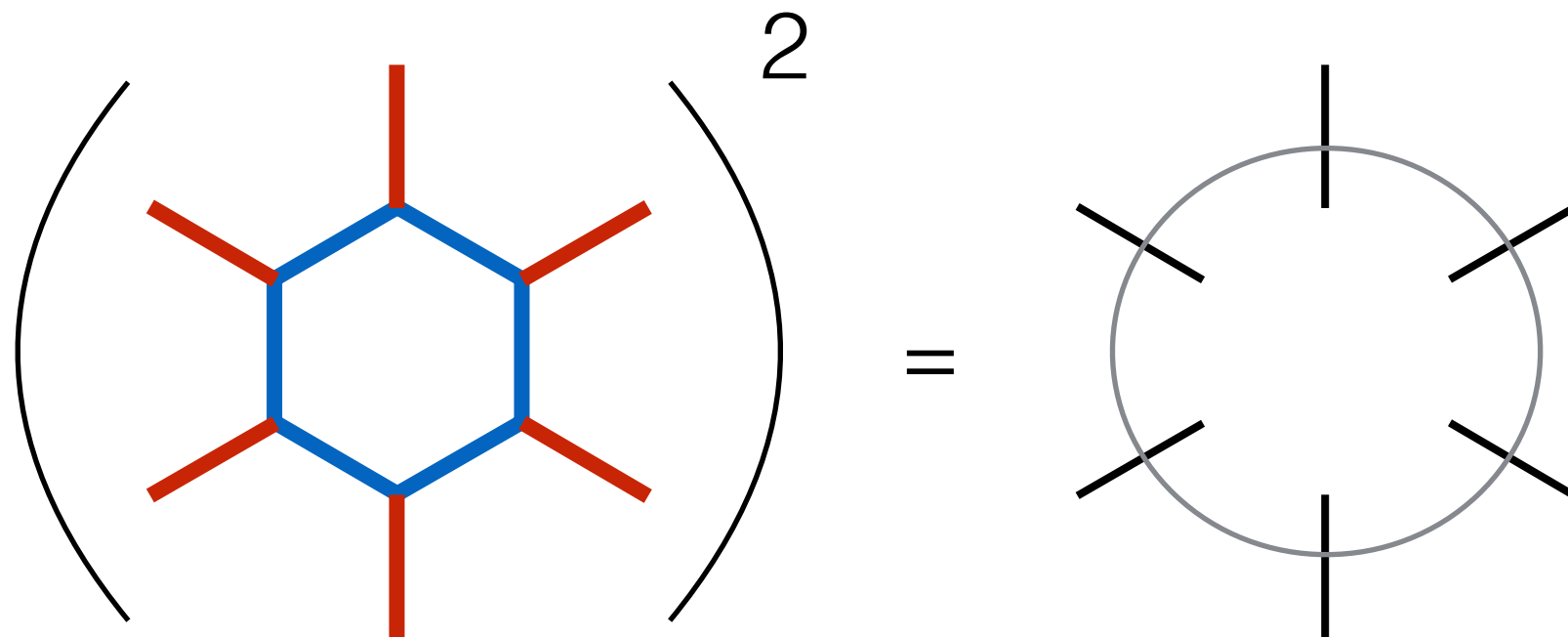
**X**  
**S**  
**Z**



- The square is equal to 1 inside gauge invariant subspace

# Plaquette operator

**X**  
**S**  
**Z**



- The square is equal to 1 inside gauge invariant subspace
- Eigenvalue of original operator is  $\pm 1$  inside the subspace

# Commutator

$$[X, S] = XSX^{-1}S^{-1} = iZ$$

$X$

$S$

$Z$

$XS$

$XS^3$

# Commutator

$$[X, S] = XSX^{-1}S^{-1} = iZ$$

$$\left[ \textcolor{blue}{/} \textcolor{red}{/} \right] = \textcolor{black}{/}$$

$X$   
 $S$   
 $Z$   
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# Commutator

$$[X, S] = XSX^{-1}S^{-1} = iZ$$

$$\left[ \begin{array}{c} \text{blue} \\ \text{red} \end{array} \right] = \text{black}$$

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$X$   
 $S$   
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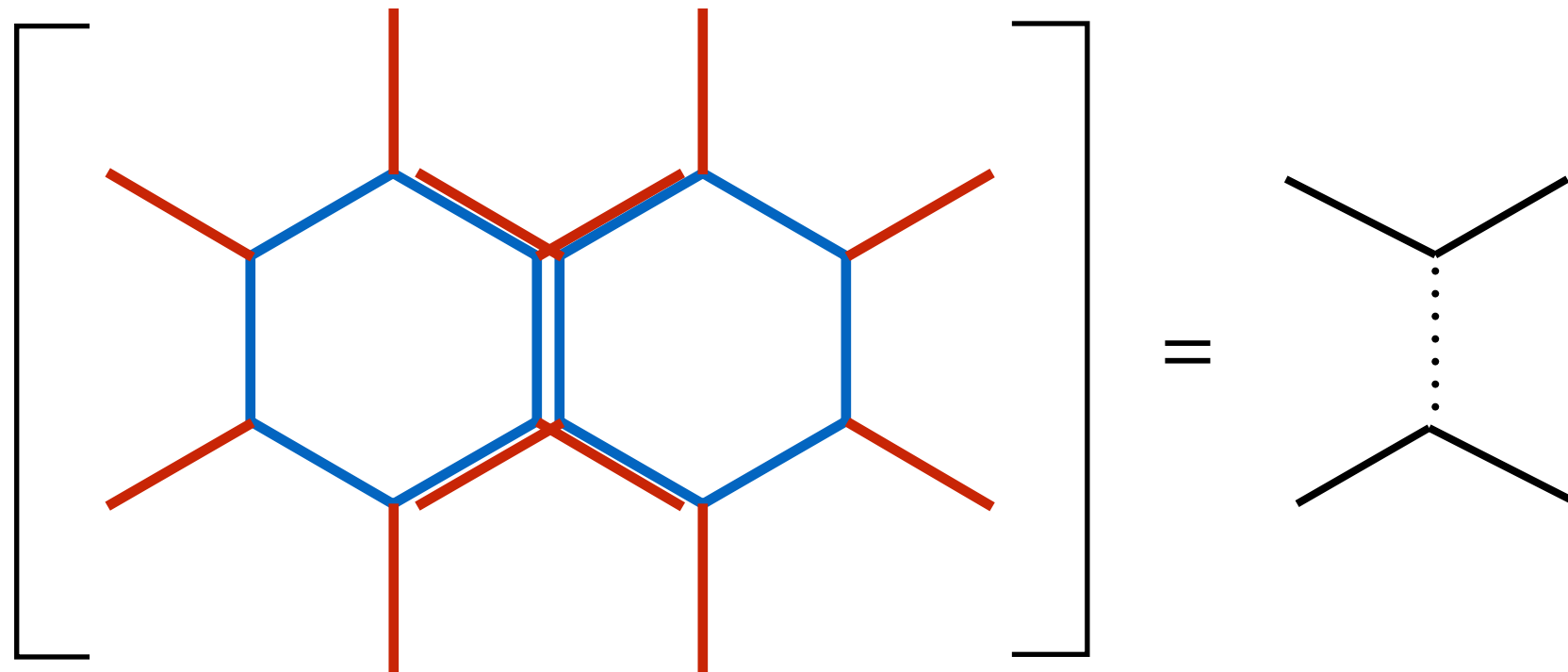
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Different non-black color: **Z**

**X**  
**S**  
**Z**  
**XS**  
**XS<sup>3</sup>**

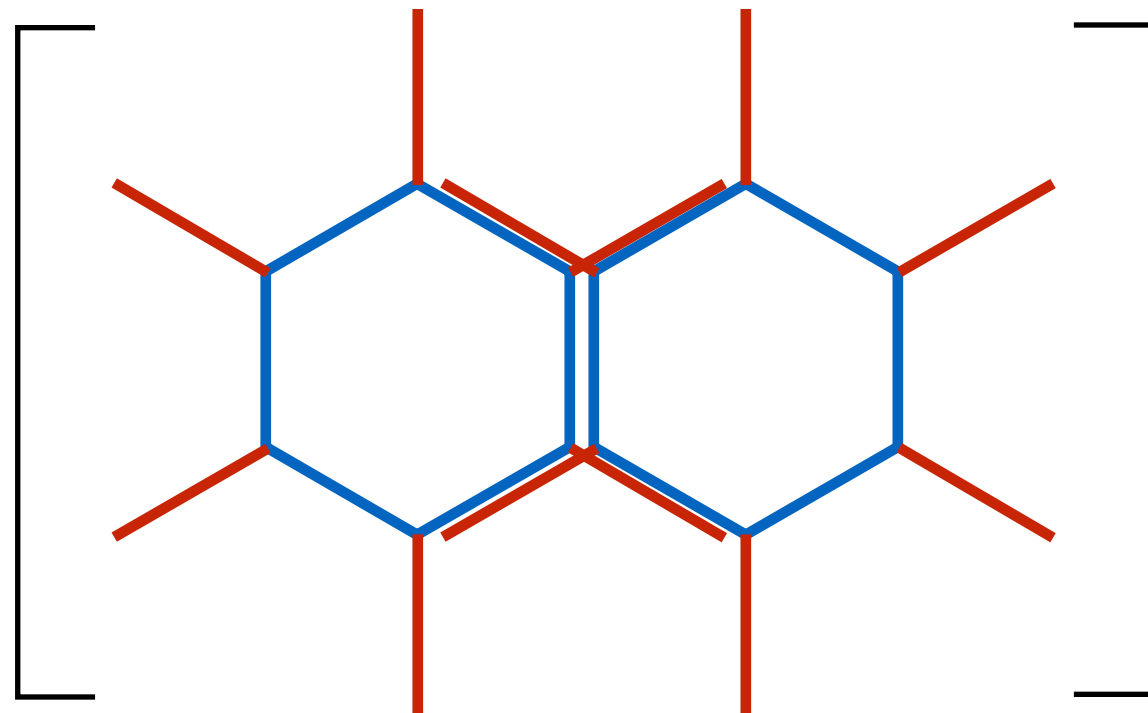
# Commutator



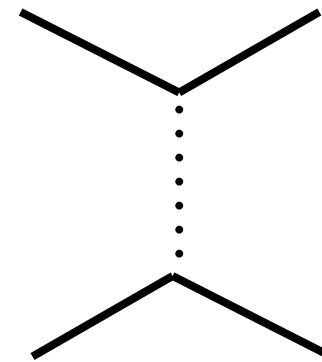
$X$   
 $S$   
 $Z$

Different non-black color:  **$Z$**

# Commutator



=



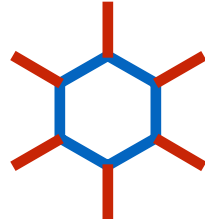
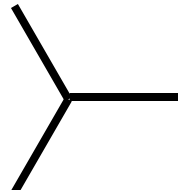
+1 inside the subspace

**x**  
**s**  
**z**

Different non-black color: **z**

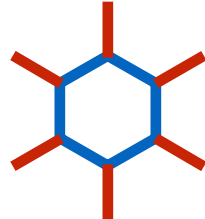
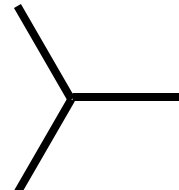
# Commuting Hamiltonians

**X**  
**S**  
**Z**

- The Plaquette operators  are hermitian and commuting in the gauge invariant subspace
- The gauge invariant subspace = locally project into the +1 eigenspace of Z-type operators 

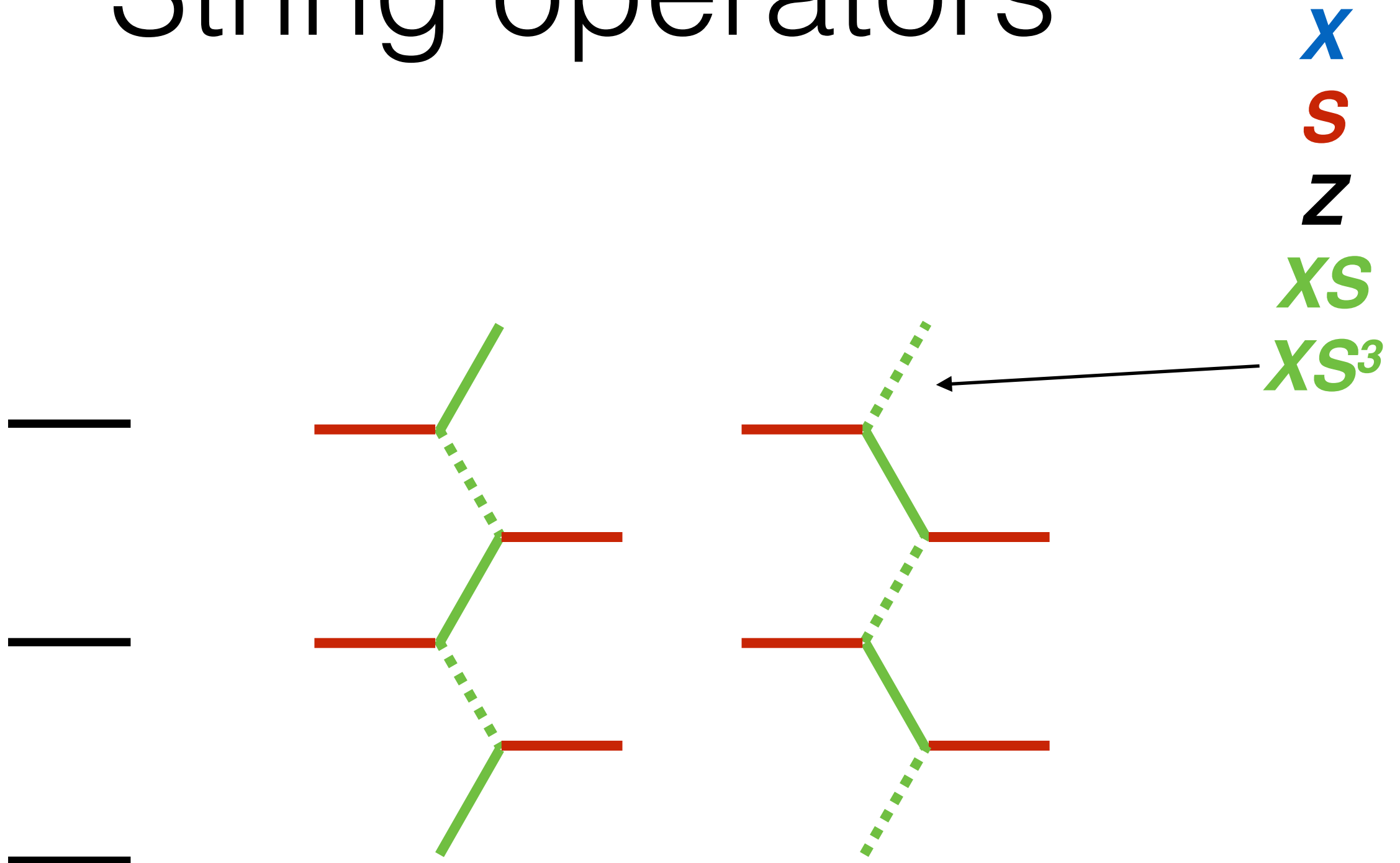
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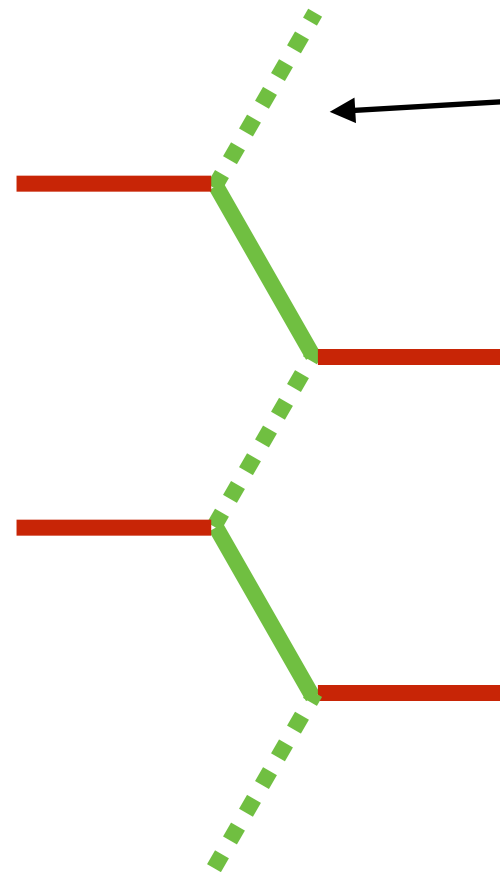
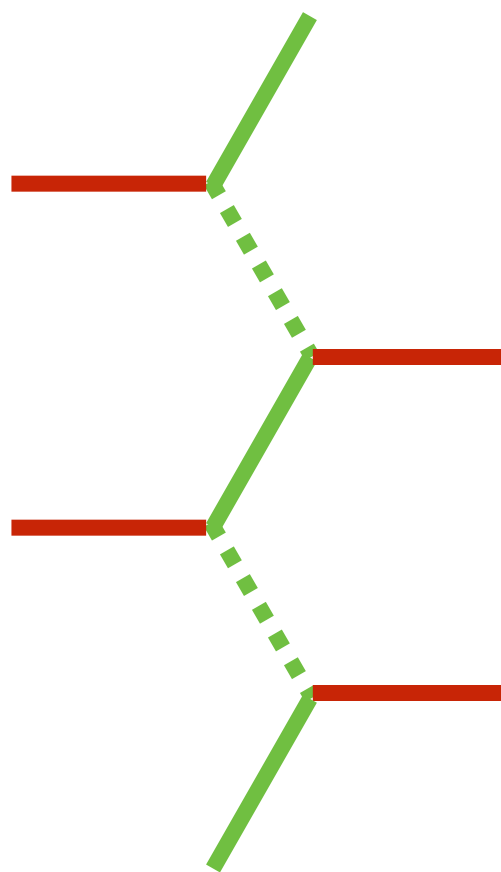
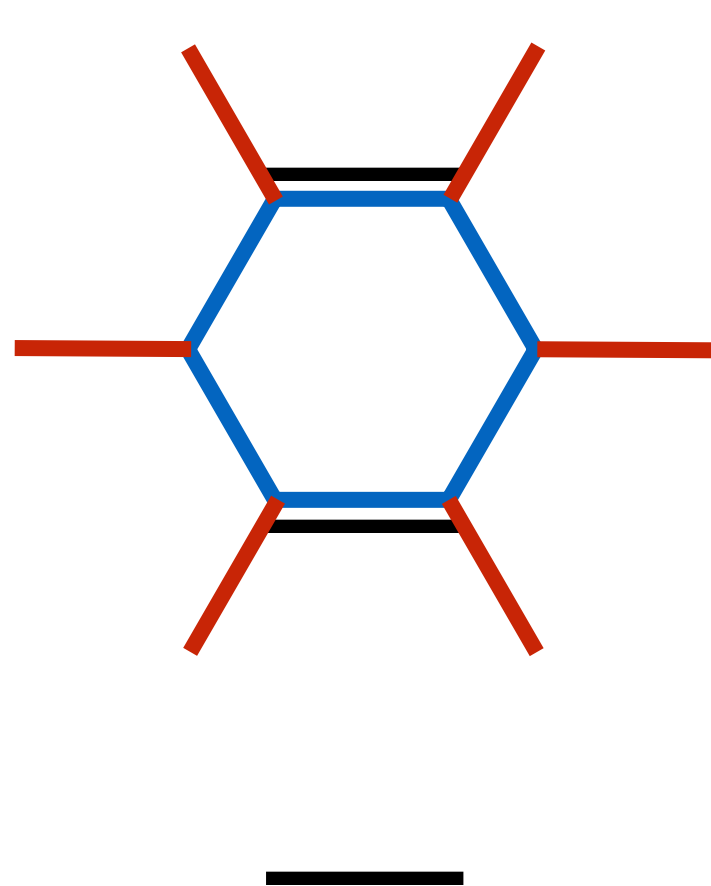
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This is a general procedure!

# String operators

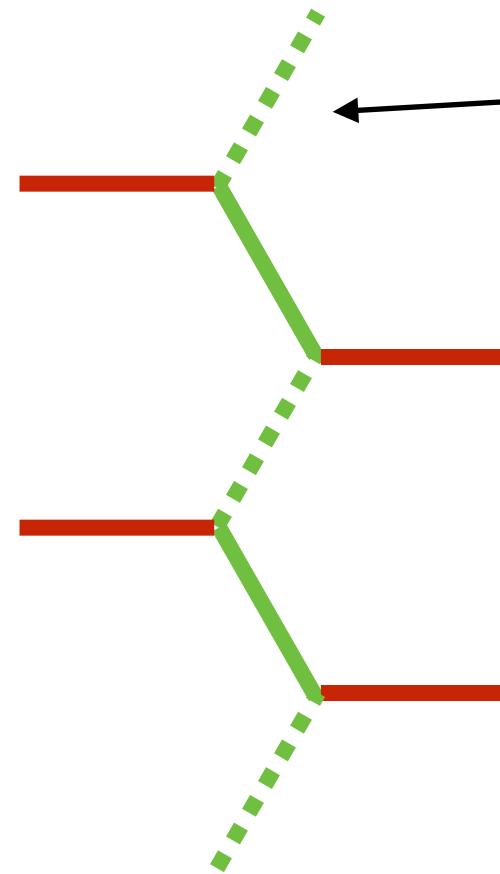
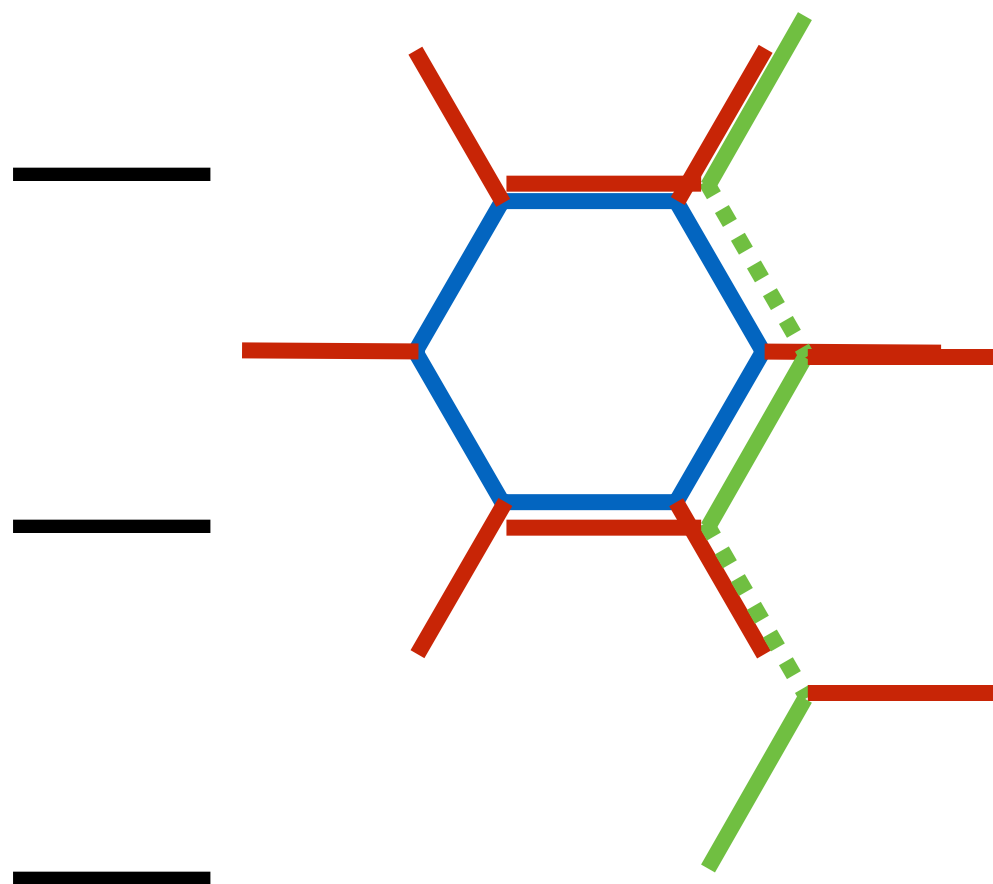


# String operators



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 $s$   
 $z$   
 $xs$   
 $xs^3$

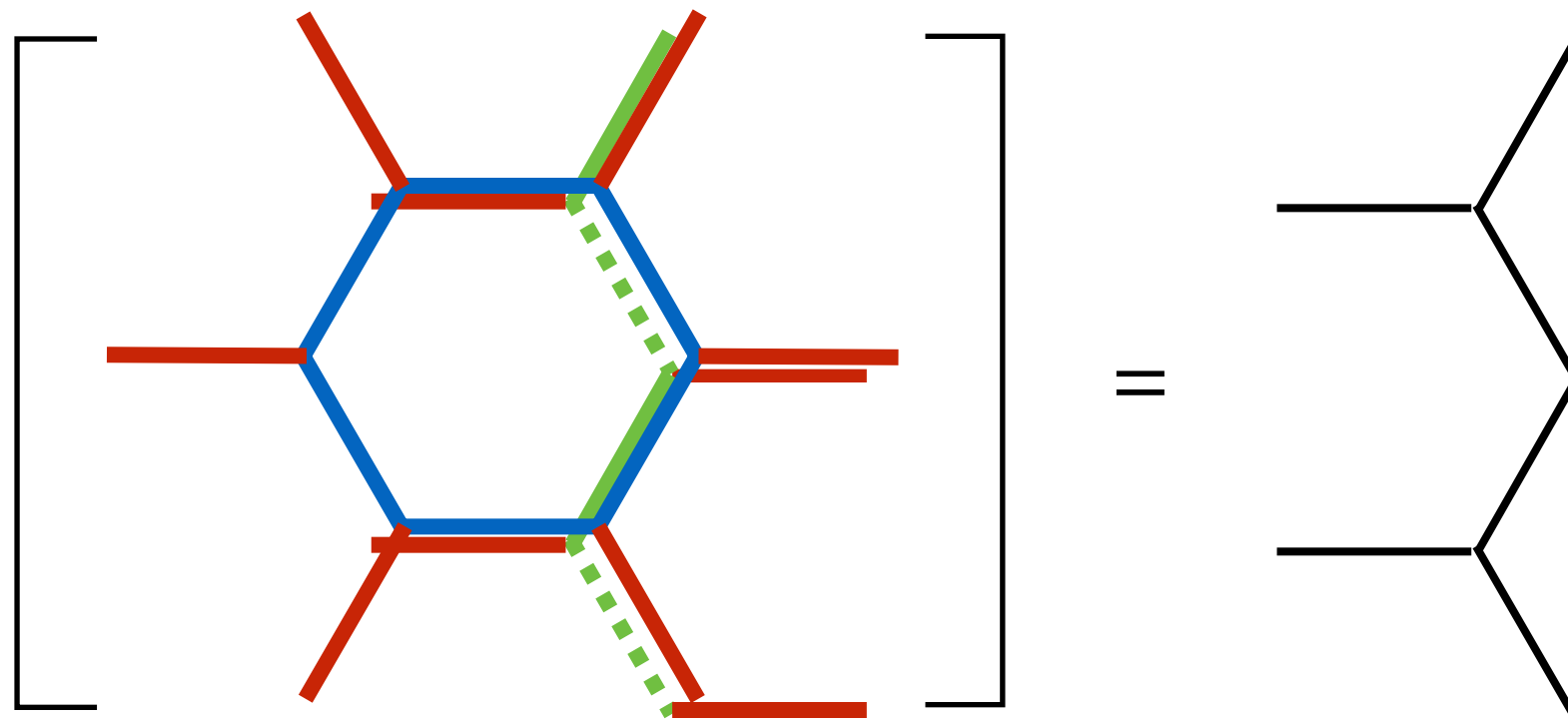
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# Commutator

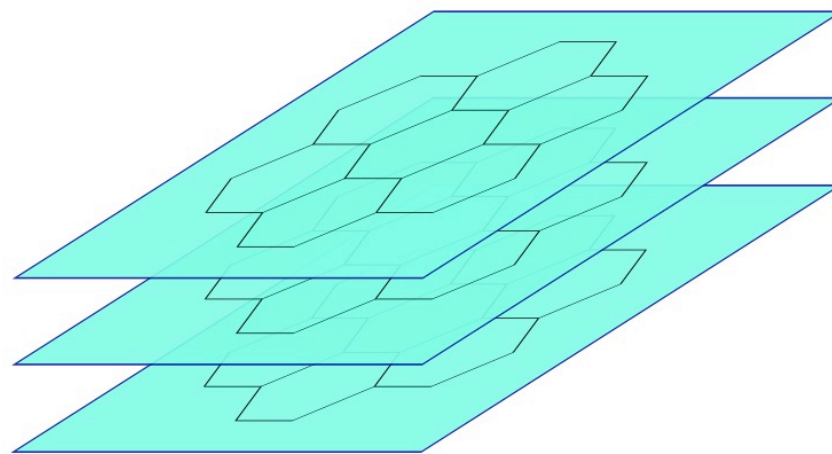


$X$   
 $S$   
 $Z$   
 $XS$   
 $XS^3$

Different non-black color:  **$Z$**

# Twisted quantum double

- Closed loops on each layer, with a phase add to each configuration
- (A subclass) can be described by XS stabilizer. Some of them support non-abelian anyons.



# Summary of properties

# Computational complexity

- Given  $G = \langle g_1, \dots, g_m \rangle \subset \mathcal{P}_n^s$ , is there a state stabilized by it?

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$$i^3 S_j \otimes S_k \otimes S_l$$

$\vdots$

NP-complete

1 in 3 SAT

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NP-complete  
1 in 3 SAT

Diagonal stabilizers  
have no **S**

Efficient

Degeneracy  $2^k$

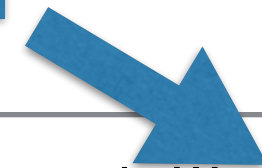
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Double semion



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# Form of the state

- We can construct a specific basis  $\{\psi_j\}$  for the code space.
- For each  $\psi_j$ , we can efficiently find a circuit of (first) Clifford and (then)  $\{T, CS, CCZ\}$  which generate the state
- $\langle \psi_j | P | \psi_k \rangle$  can be computed for Pauli operator  $P$  efficiently.



# Entanglement property

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- For a given XS-stabilizer state  $|\psi\rangle$  and a bipartition  $(A, B)$ , we can efficiently find a Pauli state  $|\varphi_{AB}\rangle$  and  $U_A \otimes U_B$  such that  $U_A \otimes U_B |\psi\rangle = |\varphi_{AB}\rangle$ .

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- Indirect way to compute entropy for XS states.
- Reflects the fact that toric code and double semion have very similar entanglement properties.

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- A more generalized (and interesting) stabilizer formalism?

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- A more generalized (and interesting) stabilizer formalism?
- A larger class of commuting projector problems that in **NP**
- Understanding entanglement properties better

Thanks



Sydney modern art museum